# Admissible vertices of contraction-critically 5-connected graphs

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# abstract

Let G be a 5-connected graph. An edge of a G is said to be 5-contractible if the contraction of the edge results in a 5-connected graph. If G has no 5-contractible edge, then it is said to be contraction-critical. An induced subgraph A of G is said to be a fragment if |N(A)| = 5 and  $V(G) - (A \cup N(A)) \neq \emptyset$ , where N(A) is the neighborhood of A. For a fragment A and  $x \in N(A)$ , a vertex  $z \in N(x) \cap N(A)$  is said to be an admissible vertex for (x; A), if the degree of z is 5 and  $|N(z) \cap A| \ge 2$ . We show some new properties on admissible vertices of contraction-critically 5-connected graphs. Using admissible vertices, we give a result on the structure around a fragment whose cardinality is 2.

Key Words: 5-connected graph, contraction-critically 5-connected, degree 5 vertex AMS classification: 05C40

Dedicated to Professor Hideo Osawa on the occasion of his retirement.

### 1 Introduction

In this paper, we deal with finite undirected graphs with neither self-loop nor multiple edge. For a graph G, let V(G) and E(G) denote the set of vertices of G and the set of edges of G, respectively. We call |V(G)| and |E(G)| the order of G and the size of G, respectively. Let  $V_k(G)$  denote the set of vertices of degree k. For an edge  $e \in E(G)$ , we denote the set of end vertices of e by V(e). For a vertex  $x \in V(G)$ , we denote by  $N_G(x)$  the neighborhood of x in G. Moreover, for a subset  $S \subseteq V(G)$ , let  $N_G(S) = \bigcup_{x \in S} N(x) - S$ . We denote the degree of  $x \in V(G)$  by  $\deg_G(x)$ . For a vertex  $x \in V(G)$ , we denote by  $E_G(x)$  the set of edges incident with x. Then  $\deg_G(x) = |N(x)| = |E_G(x)|$ . When there is no ambiguity, we write  $V_k$ , N(x), N(S),  $\deg(x)$  and E(x) for  $V_k(G)$ ,  $N_G(x)$ ,  $N_G(S)$ ,  $\deg_G(x)$  and  $E_G(x)$ , respectively. For  $S \subseteq V(G)$ , we let G[S] denote the subgraph induced by S in G. For  $S \subseteq V(G)$ , we let G - S denote the graph obtained from G by deleting the vertices in S together with the edges incident with them; thus G - S= G[V(G) - S]. A subset  $S \subseteq V(G)$  is said to be a *cutset* of G, if G - S is not connected. A cutset S is said to be a *k*-*cutset* if |S| = k. For a noncomplete connected graph G, the order of a minimum cutset of G is said to be the connectivity of G denoted by  $\kappa(G)$ . Let G be a connected graph with  $\kappa(G) = k$ . We denote by  $K_n$  the complete graph on n vertices. For graphs G and H, we write G + H the join of G and H.

Let k be an integer such that  $k \ge 2$  and let G be a k-connected graph with  $|V(G)| \ge k+2$ . An edge e of G is said to be k-contractible if the contraction of the edge results in a k-connected graph. Note that, in the contraction, we replace each resulting pair of double edges by a simple edge. If an edge is not k-contractible, then it is called k-noncontractible. Note that an edge e of G is k-noncontractible if and only if there is a k-cutset S of G such that  $V(e) \subseteq S$ . If G does not have a k-contractible edge, then G is said to be contraction-critically k-connected.

An induced subgraph A of G is called a *fragment* if |N(A)| = k and  $V(G) - (A \cup N(A)) \neq \emptyset$ . If |A| = i, then a fragment A is called *i*-fragment. A noncontractible edge e is said to be trivial, if there is a fragment A such that |A| = 1 and  $V(e) \subseteq N(A)$ . A noncontractible edge e is said to be far from trivial, if  $|A| \ge \frac{1}{2}(|V(G)| - 2k)$  for any fragment A such that  $V(e) \subseteq N(A)$ .

Let G be a 5-connected graph. Let  $x \in V(G)$  and let A be a fragment of G such that  $x \in N(A)$ . For  $y \in N(x) \cap A$ , a vertex z is said to be an *admissible vertex for* (x, y; A), if  $z \in N(x) \cap N(y) \cap S \cap V_5$  and  $|N(z) \cap A| \ge 2$ . A vertex z is said to be an *admissible vertex for* (x; A), if z is an admissible vertex for (x, y; A) for some  $y \in N(x) \cap A$ .

It is known that every 3-connected graph of order 5 or more contains a 3-contractible edge [13]. There are infinitely many contraction-critically k-connected graphs for each  $k \ge 4$  [12]. It is known that a 4-connected graph G is contarction-critical if and only if G is 4-regular, and for each edge e of it, there is a triangle which contains e [8, 10].

Egawa determined the following sharp minimum degree condition for a *k*-connected graph to have a *k*-contractible edge.

**Theorem A** (Egawa [7]) Let k be an integer, let G be a k-connected graph with  $\delta(G) \ge \left\lfloor \frac{5n}{2} \right\rfloor$ . Then G has a k-contractible edge, unless  $2 \le k \le 3$  and G is isomorphic to  $K_{k+1}$ .

There are infinitely many contraction-critically 5-connected graphs which are not 5-regular. However, by virtue of Theorem A, we know that the minimum degree of a contraction-critically 5-connected graph is 5.

The following result due to Su says that there are degree 5 vertices everywhere in a

contraction-critically 5-connected graph.

**Theorem B** (Su [11]) Every vertex of a contraction-critically 5-connected graph has two neighbors of degree five.

Since a contraction-critically 4-connected graph is 4-regular, it has very restricted substructure. On the other hand, for any given graph, there is a contraction-critically 5-connected graph which has it as an induced subgraph.

**Theorem C** (Ando and Kawarabayashi [6]) Let k be an integer such that  $k \ge 5$  and let H be a graph. Then, we can construct a contraction-critically k-connected graph which contains H as an induced subgraph.

Theorem C indicates the big difference between 'contraction-critically 4-connected graphs' and 'contraction-critically 5-connected graphs'. As Kriesell wrote in [9], it is probably a tremendously hard problem to characterize contraction-critically k-connected graphs for  $k \ge 5$ . Although we still do not have enough knowledge of the global structure of contraction-critically 5-connected graphs, we have a local structure theorem on contraction-critically 5-connected graphs [1] and we also have some progress on the study of contraction-critically 5-connected graphs [3, 4, 5, 6]. In the last decade, in the study of contraction-critically 5-connected graphs, 'admissible vertices' play crucial roles. In this paper we focus on admissible vertices of contraction-critically 5-connected graphs and we show some new conditions for a contraction-critically 5-connected graph to have an admissible vertex. Furthermore, using admissible vertices, we prove the following Theorem 1 which shows the remarkable structure around a connected 2-fragment of a contraction-critically 5-connected graph.

**Theorem 1** Let G be a contraction-critically 5-connected graph. Let A be a connected fragment of G with |A| = 2, say  $A = \{x_1, x_2\}$  and let S = N(A).

(1) If  $A \cap V_6 \neq \emptyset$ , then the number of vertices  $y \in S$  such that there is an admissible vertex for (y; A) is greater than or equal to 4.

(2) If  $A \cap V_6 \neq \emptyset$ , then the number of admissible vertices for some (y; A) is greater than or equal to 3.

(3) If  $A \cap V_6 = \emptyset$ , then there is a vertex  $y \in S - N(x_1) \cap N(x_2)$  such that there is an admissible vertex for (y; A).

This paper consists of 4 sections. After presenting preliminary results in section 2, we give some sufficient conditions for the existence of admissible vertices for given pair (x, A), where A is a fragment of a contraction-critically 5-connected graph and  $x \in N(A)$ , in section 3. In section 4, we give a proof of Theorem 1.

To conclude the section, we present three contraction-critically 5-connected graphs. The first one is 5-regular, and for each edge e of it, there is a triangle which contains e. Hence, this graph is similar in structure to contraction-critically 4-connected graphs. The second one has large maximum degree. The last one has an edge which is far from trivial. We observe that every edge in a contraction-critically 4-connected graph is trivial and, every edge of the first example and the second example is trivial. However the number of non-trivial noncontractible edges of the last example is proportional to the size of it.

### Example 1

Identifying the top and the bottom, and the left side and the right side of the graph in Fig 1, we obtain a 5-regular contraction-critically 5-connected graph for each edge e of which, there is a triangle containing e.



Fig.1: A contraction-critically 5-connected graph similar in structure to contraction contractioncritically 4-connected graphs

# Example 2

Let *H* be a contraction-critically 4-connected graph and let  $G = H + K_1$ . Then, we observe that *G* is 5-connected and every edge of *G* is trivially 5-noncontractible. Hence *G* is a contraction-

critically 5-connected graph with  $\Delta(G) = |V(G)| - 1$ .

### Example 3

Let  $K_4^-$  stand for the graph obtained from  $K_4$  by removing one edge; that is  $K_4^- \cong K_2^- + 2K_1^-$ . Let m be an integer such that  $m \ge 3$  and we construct  $G^{(m)}$  as follows; At first we prepare a configuration  $H_m$  which consists of m copies of  $K_4^-$  (see in Fig.2). Next take other three distinct vertices and join them to bottom part vertices of  $H_m$ . At last take two distinct  $K_4^-$ 's and join one  $K_4^-$  to the left side 2 vertices of  $H_m$  and the three distinct vertices, and join the other  $K_4^-$  to the right side 2 vertices of  $H_m$  and the three distinct vertices, appropriately (see Fig.3).



We call the resulting graph  $G^{(m)}$ . Let e be an edge of the top part of  $H_m$ . Then we observe that there is a 5-cutset of  $G^{(m)}$  consisting of V(e) and the distinct three vertices. Moreover, we observe that this is the only 5-cutset in  $G^{(m)}$  which contains V(e). By these observations, we know that  $G^{(m)}$  is a contraction-critically 5-connected graph and it has a far from trivial edge and many non-trivial 5-noncontractible edges.



Fig.3:  $G^{(6)}$  : contraction-critically 5-connected graph with many non-trivial edges

# 2 Preliminaries

In this section we give some more definitions and preliminary results.

For a graph G, we denote |G| for |V(G)|. For a subgraphs A and B of a graph G, when there is no ambiguity, we write simply A for V(A) and B for V(B). So N(A) and  $A \cap B$  mean N(V(A))and  $V(A) \cap V(B)$ , respectively. Also for a subgraph A of G and a subset S of V(G) we write  $A \cap S$ and  $A \cup S$  for  $V(A) \cap S$  and  $V(A) \cup S$ , respectively. When there is no ambiguity, we write E(S)for E(G[S]). For subset S and T of V(G), we denote the set of edges between S and T by  $E_G(S,T)$ . We write  $E_G(x,T)$  for  $E_G(\{xT\})$ . When there is no ambiguity, we write E(S,T) and E(x,T) for  $E_G(S,T)$  and  $E_G(x,T)$ , respectively. Let  $V_{\geq k}(G)$  (or sometimes simply  $V_{\geq k}$ ) denote the set of vertices of degree at least k.

Let G be a connected graph with  $\kappa(G) = k$ . Recall that an induced subgraph A of G is called a fragment if |N(A)| = k and  $V(G) - (A \cup N(A)) \neq \emptyset$ . In other words, a fragment A is a nonempty union of components of G - S where S is a k-cutset of G such that  $V(G) - (A \cup S) \neq \emptyset$ . By the definition, if A is a fragment of G, then  $G - (A \cup N(A))$  is also a fragment of G. Let  $\overline{A}$  stand for  $G - (A \cup N(A))$ . For an edge e of G, a fragment A of G is said to be a *fragment with respect* to e if  $V(e) \subseteq N(A)$ . For a set of edges  $F \subseteq E(G)$ , we say that A is a *fragment with respect to* F if A is a fragment with respect to some  $e \in F$ . A fragment A with respect to F is said to be *minimum* (resp. *minimal*) if there is no fragment B other than A with respect to F such that |B| < |A| (resp.  $B \subsetneq A$ ). If |A| = 1, then a fragment A is said to be *trivial*.

Let  $V_k^{(i)}(G)$  (or sometimes simply  $V_k^{(i)}$ ) stand for the set of vertices of  $V_k(G)$  each of which has i neighbors in  $V_k(G)$ , namely  $V_k^{(i)} = \{x \in V_k(G) \mid | N(x) \cap V_k(G) \mid = i\}$ .

We start with the following Lemma 1 [3] which is a simple but useful observation. We give a proof of Lemma 1 for the reader's convenience.

**Lemma 1** Let A be a fragment of a k-connected graph G and let  $S \subseteq N(A)$ . If  $|N(S) \cap A| < |S|$ , then  $A = N(S) \cap A$ .

*Proof.* Assume that  $A \neq N(S) \cap A$ . Let  $A' = A - (N(S) \cap A)$ . Since  $A' \neq \emptyset$  and  $T \cap (\overline{A} \cup S) = \emptyset$ ,  $(N(A) - S) \cup (N(S) \cap A)$  separates A' and  $\overline{A} \cup S$ . Since  $|N(S) \cap A| < |S|$ , we see that  $|(N(A) - S) \cup (N(S) \cap A) = |N(A)| - |S| + |N(S) \cap A| < |N(A)| = k$ , which contradicts the *k*-connectedness of *G*.

The reader can find the proof of Lemma 2 in [3].

**Lemma 2** Let G be a 5-connected graph, and let A and B be fragments of G. Let S = N(A) and let T = N(B).

В	$\bar{A} \cap B$	$S \cap B$	$A \cap B$
T	$\bar{A} \cap T$	$S \cap T$	$A \cap T$
$\bar{B}$	$\bar{A} \cap \bar{B}$	$S\cap \bar{B}$	$A \cap \bar{B}$
	Ā	S	A

# Then the following hold.

(1) If  $|(S \cap B) \cup (S \cap T) \cup (A \cap T)| \ge 6$ , then  $|(\bar{A} \cap T) \cup (S \cap T) \cup (S \cap \bar{B})| \le 4$  and  $\bar{A} \cap \bar{B} = \emptyset$ . In particular, if neither  $A \cap B$  nor  $\bar{A} \cap \bar{B}$  is empty, then both  $A \cap B$  and  $\bar{A} \cap \bar{B}$  are fragments of G.

(2)  $|(S \cap B) \cup (S \cap T) \cup (A \cap T)| = 5 + |S \cap B| - |\overline{A} \cap T|$ . In particular, if  $A \cap B \neq \emptyset$ , then  $|S \cap B| \ge |\overline{A} \cap T|$ .

(3) If  $|\bar{A}| \ge 2$ , then either  $|(S \cap B) \cup (S \cap T) \cup (A \cap T)| \le 5$  or  $|(S \cap \bar{B}) \cup (S \cap T) \cup (A \cap T)| \le 5$ .

# 3 Admissible vertices

In the following two sections we consider 5-connected graphs.

We introduce 'admissible vertex' in [3] and we introduce 'strongly admissible vertex' and 'hyper admissible vertex' in [2]. In this paper, we introduce 'insufficient' and give a new sufficient condition a contraction-critically 5-connected graph to have an admissible vertex.

Let G be a 5-connected graph. Let  $x \in V(G)$  and let A be a fragment of G such that  $x \in N(A)$ . Let S = N(A).

Let  $y \in N(x) \cap A$ . Recall that a vertex z is said to be an admissible vertex for (x, y; A), if the following two conditions hold.

(1)  $z \in N(x) \cap N(y) \cap S \cap V_5$ .

(2)  $|N(z) \cap A| \ge 2$ .

Here, we introduce more detailed properties of admissible vertices.

For  $y \in N(x) \cap A$ , a vertex z is said to be an *strongly admissible vertex for* (x, y; A), if the following conditions hold.

- (1)  $z \in N(x) \cap N(y) \cap S \cap V_5$ ,
- (2)  $|N(z) \cap A| \ge 2$ , and
- (3)  $|N(z) \cap \overline{A}| = 1$ .

For  $y \in N(x) \cap A$ , a vertex z is said to be an hyper admissible vertex for (x, y; A), if the following conditions hold.

- (1)  $z \in N(x) \cap N(y) \cap S \cap V_5$ ,
- (2)  $|N(z) \cap A| \ge 2$ , and
- (3)  $|N(z) \cap \overline{A}| = |N(z) \cap S| = 1$ .

A vertex z is said to be a strongly admissible vertex for (x; A) or a hyper admissible vertex for (x; A), if z is a strongly admissible vertex for (x, y; A) or a hyper admissible vertex for (x, y; A)for some  $y \in N(x) \cap A$ , respectively.

A triangle H of G is said to be an A-inner  $x^*$ -triangle if (1)  $x \in V(H)$ , (2)  $V(H) - \{x\} \subseteq A$ and (3)  $(V(H) - \{x\}) \cap V_5 \neq \emptyset$ .

A vertex x is said to be *insufficient on A* if the following two conditions hold.

(1) there is no A-inner  $x^*$  -triangle.

(2)  $N(u) \cap N(u') \cap A \cap V_5 = \emptyset$  for any  $u, u' \in N(x) \cap A \cap V_5$ .

The following Lemmas 3 and 4 give some basic properties of admissible vertices in a contraction-critically 5-connected graph. The reader can find proofs of Lemmas 3 and 4, and Corollary 7 in [3], however for the convenience of the reader, we give proofs of Lemmas 3 and 4. We give an alternate proof of Corollary 7 in this section.

**Lemma 3** ([3] Corollary 4) Let x be a vertex of a contraction-critically 5-connected graph G. Let A be a fragment of G such that  $x \in N(A)$ . Suppose  $|\bar{A}| \ge 2$ ,  $|A| \ge 3$  and  $|N(x) \cap A| = 1$ . Then, there is an admissible vertex for (x; A)

*Proof.* Let  $N(x) \cap A = \{y\}$ . Let B be a fragment with respect to xy. Let S = N(A) and let T = N(B). Since  $|\bar{A}| \ge 2$ , by Lemma 2 (3), we see that either  $|(S \cap B) \cup (S \cap T) \cup (A \cap T)| \le 5$  or  $|(S \cap \bar{B}) \cup (S \cap T) \cup (A \cap T)| \le 5$ . Without loss of generality we may assume  $|(S \cap B) \cup (S \cap T) \cup (A \cap T)| \le 5$ . Then, since  $N(x) \cap A = \{y\}$ , we have  $A \cap B = \emptyset$ .

Claim 3.1  $|(S \cap \overline{B}) \cup (S \cap T) \cup (A \cap T)| \ge 6.$ 

*Proof.* If  $A \cap \overline{B} \neq \emptyset$ , then  $|(S \cap \overline{B}) \cup (S \cap T) \cup (A \cap T)| \ge 6$  since  $N(x) \cap A = \{y\}$ . Hence, we assume  $A \cap \overline{B} = \emptyset$ . Then, since  $A \cap B = \emptyset$ , we have  $A = A \cap T$  and  $|A| = |A \cap T| \ge 3$ , which

implies that  $|A \cap T| > |S \cap B|$  since  $|(S \cap B) \cup (S \cap T) \cup (A \cap T)| \le 5$ .

Hence we observe that  $|(S \cap \overline{B}) \cup (S \cap T) \cup (A \cap T)| > |S| = 5$  and Claim 3.1 is proved. Claim 3.1 assures us that  $|A \cap T| > |S \cap B|$ . If  $|S \cap B| \ge 2$ , then  $|A \cap T| \ge 3$  and  $|(S \cap B) \cup (S \cap T) \cup (A \cap T)| \ge 6$ , which contradicts the fact that  $|(S \cap B) \cup (S \cap T) \cup (A \cap T)| \le 5$ . Hence  $|S \cap B| \le 1$ . Claim 3.1 also assures us that  $\overline{A} \cap B = \emptyset$  and  $B = S \cap B$ . Let  $B = S \cap B = \{z\}$ . Then we observe that z is an admissible vertex for (x; A).

**Lemma 4** ([3] Lemma 3) Let x be a vertex of a contraction-critically 5-connected graph G. Let A be a fragment such that  $x \in N(A)$ ,  $|\overline{A}| \ge 2$  and  $|A| \ge 3$ . Then, for each vertex  $y \in N(x) \cap A$ , there is either an admissible vertex for (x, y; A) or a fragment A' with respect to xy such that  $A' \subsetneq A$ .

*Proof.* Assume that there is neither an admissible vertex for (x, y; A) nor a fragment A' with respect to xy such that  $A' \subsetneq A$ . Let B be a fragment with respect to xy. Let S = N(A) and let T = N(B). Since  $|\bar{A}| \ge 2$ , by Lemma 2 (3), we see that either  $|(S \cap B) \cup (S \cap T) \cup (A \cap T)| \le 5$  or  $|(S \cap \bar{B}) \cup (S \cap T) \cup (A \cap T)| \le 5$ . Without loss of generality we may assume  $|(S \cap B) \cup (S \cap T) \cup (A \cap T)| \le 6$ . Without loss of generality we may assume  $|(S \cap B) \cup (S \cap T) \cup (A \cap T)| \le 6$ . If  $A \cap B \ne \emptyset$ , then  $A \cap B$  is a fragment with respect to xy such that  $A \cap B \subsetneq A$  since  $y \in A \cap T$ , which contradicts the assumption. Hence  $A \cap B = \emptyset$ .

Claim 4.1  $A \cap \overline{B} \neq \emptyset$ .

*Proof.* Assume  $A \cap \overline{B} = \emptyset$ . Then  $A = A \cap T$  and  $|A \cap T| = |A| \ge 3$ . Hence  $|S \cap B| = |(S \cap B) \cup (S \cap T) \cup (A \cap T)| - |S \cap T| - |A \cap T| \le 5 - 1 - 3 = 1$ . Thus  $|S \cap B| = 1$ , say  $S \cap B = \{z\}$ . Then, we find that z is an admissible vertex for (x, y; A), which contradicts the assumption.

By Claim 4.1, we know that  $A \cap \overline{B} \neq \emptyset$ . Hence, if  $|(S \cap \overline{B}) \cup (S \cap T) \cup (A \cap T)| = 5$ , then  $A \cap \overline{B}$  a fragment with respect to xy such that  $A \cap \overline{B} \subsetneq A$ , which contradicts the assumption. Thus we have  $|(S \cap \overline{B}) \cup (S \cap T) \cup (A \cap T)| \ge 6$ , which implies  $\overline{A} \cap B = \emptyset$  and  $|S \cap B| < |A \cap T|$ . Therefore,  $B = S \cap B$  and  $|S \cap B| = |(S \cap B) \cup (S \cap T) \cup (A \cap T)| - |S \cap T| - |A \cap T| \le 4 - |A \cap T| < 4 - |S \cap B|$ . Hence we have  $|B| = |S \cap B| = 1$ , say  $B = S \cap B = \{z\}$ . Then, we again find that z is an admissible vertex for (x, y; A), which contradicts the assumption. This contradiction proves Lemma 4.

**Lemma 5** Let x be a vertex of a contraction-critically 5-connected graph G. Let A be a fragment with respect to E(x) such that  $|\overline{A}| \ge 2$ , |A| = 2. If there is neither an A-inner  $x^*$  -triangle nor an admissible vertex for (x; A), then,  $A \subseteq V_5$ .

*Proof.* Let  $A = \{u, u'\}$  and assume that either  $u \notin V_5$  or  $u' \notin V_5$ . Let  $S = N(A) = \{x, w, w', w'', w'''\}$ . We may assume  $u \in N(x) \cap A$ . Since A is a fragment with respect to E(x), we also assume that  $w \in N(x) \cap S$ .

Claim 5.1  $S - \{x\} \subseteq N(u')$ .

**Proof.** If  $u' \notin V_5$ , then  $N(u') = S \cup \{u\}$  and we are done. Hence assume  $u' \in V_5$ . If  $u'x \in E(G)$ , then we see that  $G[\{x, u, u'\}]$  is an A-inner  $x^*$  -triangle, which violates the assumption. Hence  $u'x \notin E(G)$ , which implies the desired conclusion,  $S - \{x\} \subseteq N(u')$ .

Let *B* be a fragment with respect to xu and let T = N(B).

**Claim 5.2** (1)  $u' \in T$  and (2)  $|S \cap B| = |S \cap \overline{B}| = 2$ .

*Proof.* (1) By Claim 5.1, we see that  $S - \{x\} \subseteq N(u')$ , which implies  $u' \in T$ .

(2) Assume  $|S \cap B| \le 1$ . Then  $\overline{A} \cap B = \emptyset$  since  $|S \cap B| < |A \cap T|$ . If  $S \cap B = \emptyset$ , then  $B = \emptyset$ , which contradicts the choice of B. Hence assume  $|S \cap B| = 1$  and let  $S \cap B = \{y\}$ . Then we see that y is an admissible vertex for (x; A), which contradicts the assumption. Hence  $|S \cap B| \ge 2$ . Similarly we see  $|S \cap \overline{B}| \ge 2$ . Then, since  $S \cap T \neq \emptyset$ , we have  $|S \cap B| = |S \cap \overline{B}| = 2$ .

By Claim 5.2 (2), we may assume that  $S \cap B = \{w, w'\}$  and  $S \cap \overline{B} = \{w'', w'''\}$ .

Claim 5.3 If  $uw \in E(G)$ , then  $w \notin V_5$ .

*Proof.* Assume that  $uw \in E(G)$  and  $w \in V_5$ . Then, by Claim 5.1, we see that  $u'w \in E(G)$ . This implies w is an admissible vertex for (x; A), which contradicts the assumption.

Claim 5.4  $u \in V_5$ .

**Proof.** Assume  $u \notin V_5$ . Then  $N(u) = S \cup \{u'\}$ . Hence  $uw \in E(G)$  and Claim 4.3 assures us that  $w \notin V_5$ . By Claim 5.1, we know that  $u'w' \in E(G)$ . Let C be a fragment with respect to u'w' and let R = N(C). Then, since  $S \subseteq N(u)$ , we see that  $u \in R$ , which implies  $\{u, u'\} \subseteq T \cap R$ .

Subclaim 5.4.1  $w \in R$ .

**Proof.** Assume  $w \notin R$ . Without loss of generality we may assume that  $w \in C$ . Then, since  $xw \in E(G)$ , we observe that  $x \in R \cup C$ . Since  $S \cap \overline{C} \neq \emptyset$  we see that  $\{w'', w'''\} \cap \overline{C} \neq \emptyset$ , which implies  $(\overline{B} \cap \overline{C}) \cap \{w'', w'''\} \neq \emptyset$  since  $\{w'', w'''\} \subseteq \overline{B}$ . Now we observe that  $w \in B \cap C$  and

 $(\overline{B} \cap \overline{C}) \cap \{w'', w'''\} \neq \emptyset$ , which implies that  $|(R \cap B) \cup (R \cap T) \cup (C \cap T)| = 5$ . Hence  $B \cap C$  is a fragment of G. Since  $\{w'', w'''\} \subseteq \overline{B}$ ,  $x \in T$  and  $w' \in R$ , we see that  $N(\{u, u'\}) \cap (B \cap C) = \{w\}$ . Hence, applying Lemma 1 with the roles of A and S replaced by  $B \cap C$  and  $\{u, u'\}$ , respectively, we see that  $C \cap B = \{w\}$ . This implies  $w \in V_5$ , which contradicts Claim 5.3. This contradiction proves Subclaim 5.4.1.

# **Subclaim 5.4.2** (1) $x \in V_5$ , and (2) $xu', xw' \in E(G)$ .

*Proof.* (1) By Subclaim 5.4.1, we know that  $\{w, w'\} \subseteq S \cap R$ , which implies either  $|S \cap C| = 1$  or  $|S \cap \overline{C}| = 1$ . Without loss of generality we may assume that  $|S \cap C| = 1$ , say  $S \cap C = \{z\}$ . Then  $z \in \{x, w'', w'''\}$ . Since  $|S \cap C| < |A \cap R|$ . Lemma 2 (2) assures us that  $\overline{A} \cap C = \emptyset$ , which implies  $C = S \cap C = \{z\}$ . Hence  $z \in V_5$  and  $zw \in E(G)$ . Since  $ww'', ww''' \notin \mathcal{E}(G)$ , we see that z = x and  $x \in V_5$ .

(2) Since N(x) = R, we observe that  $xu', xw' \in E(G)$ .

# Subclaim 5.4.3 $ww' \in E(G)$ .

*Proof.* Since  $|A \cap T| = |S \cap B| = 2$ , we see that  $|(\bar{A} \cap T) \cup (S \cap T) \cup (S \cap B)| = 5$ . Let  $N(x) = \{u, u', w, w', v\}$ . Since  $N(x) \cap \bar{A} \neq \emptyset$  and  $\{u, u', w, w'\} \subseteq (A \cap T) \cup (S \cap B)$ , we observe that  $v \in \bar{A} \cap \bar{B}$ , which implies  $N(x) \cap (\bar{A} \cap B) = \emptyset$ . Since  $|(\bar{A} \cap T) \cup (S \cap T) \cup (S \cap B)| = 5$  and  $N(x) \cap (\bar{A} \cap B) = \emptyset$ , we see that  $\bar{A} \cap B = \emptyset$ , which implies  $B = S \cap B = \{w, w'\}$ . Since  $w \notin V_5$  and  $B = \{w, w'\}$ , we have  $ww' \in E(G)$ .

We proceed with the proof of Claim 5.4. Now we observe that  $G[N(x) - \{v\}] \cong K_4$ , which implies xv is contractible. This contradicts that G is contraction-critically 5-connected and Claim 5.4 is proved.

By Claim 5.4, we have  $u \in V_5$ . Hence  $u' \notin V_5$ . But, in this situation, we see that  $G[\{x, u, u'\}]$  is an A-inner  $x^*$  -triangle, which contradicts the assumption. This contradiction proves Lemma 5.

Recall that an vertex x is said to be insufficient on a fragment A if (1) there is no A-inner  $x^*$ -triangle and (2)  $N(u) \cap N(u') \cap A \cap V_5 = \emptyset$  for any  $u, u' \in N(x) \cap A \cap V_5$ .

The following Lemma 6 says that "x is insufficient on A" is an sufficient condition for the existence of an admissible vertex for (x; A).

Lemma 6 Let x be a vertex of a contraction-critically 5-connected graph G. Let A be a fragment

such that  $x \in N(A)$ ,  $|\bar{A}| \ge 2$  and  $|A| \ge 3$ . If x is insufficient on A, then there is an admissible vertex for (x; A).

**Proof.** We prove Lemma 6 by the induction on  $|N(x) \cap A|$ . If  $|N(x) \cap A| = 1$ , then Lemma 3 assures us that the desired conclusion holds. Assume  $|N(x) \cap A| \ge 2$  and also assume that there is no admissible vertex for (x; A). Choose  $y \in N(x) \cap A$  so that  $\deg_G(y)$  to be as small as possible. Since there is no admissible vertex for (x, y; A), Lemma 5 assures us that there is a fragment A' with respect to xy such that  $A' \subsetneq A$ .

### **Claim 6.1** |A'| = 2.

**Proof.** At first assume |A'| = 1, say  $A' = \{u\}$ . Then  $u \in V_5$ ,  $\{x, y\} \subseteq N(u)$  and  $A = \{y, u\}$ . In this situation, we observe that  $G[\{x, y, u\}]$  is an A-inner  $x^*$ -triangle, which violates the fact that x is insufficient on A.

Next assume  $|A'| \ge 3$ . Then  $|A'| \ge 3$  and  $|\bar{A}'| > |\bar{A}| \ge 2$ . Since x is insufficient on A and  $A' \subsetneq A$ , x is also insufficient on A'. Since  $y \in N(x) \cap A$  and  $y \notin N(x) \cap A'$ , we see that  $|N(x) \cap A'| < |N(x) \cap A|$ . Hence, applying the induction hypothesis to A', we see that there is an admissible vertex z for (x; A'). Since  $A' \subsetneq A$ ,  $N(A') \subseteq S \cup A$ , which implies  $z \in S \cup A$ . We show  $z \in S$ . Assume that  $z \in A$ . Since z is an admissible vertex for (x; A'), there is a vertex  $u \in N(x) \cap N(z) \cap A'$ . Then, since  $z \in A \cap V_5$  and  $u \in A' \subsetneq A$ , we observe that  $G[\{x, z, u\}]$  is an A-inner  $x^*$ -triangle, which violates the fact that x is insufficient on A. Now it is shown that  $z \in S$ , which implies that z is an admissible vertex z for (x; A). This contradicts the assumption and Claim 6.1 is proved.

By Claim 6.1 we know |A'| = 2, say  $A' = \{u, u'\}$ . We may assume that  $xu \in E(G)$ . Since  $A' \subsetneq A$  and there is no A-inner  $x^*$ -triangle, we see there is no A'-inner  $x^*$ -triangle. Assume that there is an admissible vertex z for (x; A'). Then  $z \in V_5$  and  $N(x) \cap N(z) \cap A' \neq \emptyset$ . If  $z \in A$ , then we find an A-inner  $x^*$ -triangle, which contradicts the assumption. Hence  $z \in S$  and z is an admissible vertex for (x; A), which again contradicts the assumption. It is shown that there is no admissible vertex for (x; A'). Hence, there is neither an A'-inner  $x^*$ -triangle nor an admissible vertex for (x; A'). Hence, there is neither an A'-inner  $x^*$ -triangle nor an admissible vertex for (x; A'). Hence, there is neither an A'-inner  $x^*$ -triangle nor an admissible vertex for (x; A'). Thus Lemma 4 assures us that  $u, u' \in V_5$ . Recall that we choose y so that  $\deg_G(y)$  to be as small as possible. Hence, we see that  $y \in V_5$  since  $u \in N(x) \cap A \cap V_5$ . Since there is no A-inner  $x^*$ -triangle and  $y, u \in N(x) \cap A \cap V_5$ , we see that  $yu \notin E(G)$ , which implies  $uu' \in E(G)$  since  $A' = \{u, u'\}$ . If  $xu' \in E(G)$ , then  $G[\{x, u, u'\}]$  is an A-inner  $x^*$ -triangle, which contradicts the assumption. Hence  $xu' \notin E(G)$ , which implies  $yu' \in E(G)$ . Now we observe that  $y, u \in N(x) \cap A \cap V_5$  and  $u' \in N(y) \cap N(u) \cap A \cap V_5$ , which contradicts the assumption that x is

insufficient on A. This contradiction completes the proof of Lemma 6.

We note that, in the definition of 'insufficient', the condition "(2)  $N(u) \cap N(u') \cap A \cap V_5 = \emptyset$ for any  $u, u' \in N(x) \cap A \cap V_5$ " is necessary. There is a contraction-critically 5-connected graph Gwhich has a vertex x and a fragment A such that  $x \in N(A)$ ,  $|\overline{A}| \ge 2$  and  $|A| \ge 3$  and G has neither an admissible vertex for (x; A) nor an A-inner  $x^*$ -triangle.

By the definition, if  $N(x) \cap A \cap V_5 = \emptyset$ , then x is insufficient on A. Hence, the following is an immediate corollary of Lemma 6.

**Corollary 7** ([3] Lemma 6) Let G be a contraction-critically 5-connected graph G and let A be a fragment of G with  $|\bar{A}| \ge 2$  and  $|A| \ge 3$ . Let  $x \in N(A)$ . If  $N(x) \cap A \cap V_5 = \emptyset$ , then there is an admissible vertex for (x; A).

**Lemma 8** Let x be a vertex of a contraction-critically 5-connected graph G. Let A be a fragment such that  $x \in N(A)$ ,  $|\bar{A}| \ge 2$  and  $|A| \ge 3$ . Suppose  $|N(x) \cap A| = 1$  and  $N(x) \cap A \cap V_5 = \emptyset$ . Then,

(1) there is a strongly admissible vertex z for (x; A),

(2) if  $(N(z) \cap N(A) - \{x\}) \cap (V_5 - V_5^{(2)}) = \emptyset$ , then z is a hyper admissible vertex for (x; A).

*Proof.* Let S = N(A) and let  $N(x) \cap A = \{y\}$ . Note that  $y \notin V_5$  since  $N(x) \cap A \cap V_5 = \emptyset$ . By Lemma 3, there is an admissible vertex z for (x, y; A). Let  $B = \{z\}$  and let T = N(y) = N(B).

We show (1). Assume z is not strongly admissible, that is  $|N(z) \cap \overline{A}| \ge 2$ . Then, since  $z \in V_5$ , we see that  $|N(z) \cap \overline{A}| = |N(z) \cap A| = 2$ ,  $S \cap T = \{x\}$  and  $|S \cap \overline{B}| = 3$ . Let  $A \cap T = \{y, u\}$  and let  $S \cap \overline{B} = \{w, w', w''\}$ . Furthermore, let  $A' = A - \{y\}$  and  $S' = N(A') = (S - \{x\}) \cup \{y\} = \{z, y, w, w', w''\}$ . Since  $N(x) \cap A = \{y\}$ , we observe that A' is a fragment of G such that  $|A'| = |A - \{y\}| \ge 2$  and  $|\overline{A}'| = |\overline{A} \cup \{x\}| \ge 3$ . Then, since  $N(z) \cap S' = \{y\}$  and  $y \notin V_5$ , we observe that  $N(z) \cap S' \cap V_5 = \emptyset$ , which implies that there is no admissible vertex for (z; A'). If  $|A'| \ge 3$ , then Lemma 3 assures us the existence of an admissible vertex for (z; A'), which contradicts the previous assertion. Hence we have |A'| = 2, say  $A' = \{u, u'\}$ . Then  $u' \in A' \cap \overline{B}$  and  $N(u') = \{u, y, w, w', w''\}$ . Moreover we observe that  $N(u) \subseteq \{y, z, u', w, w', w''\}$  and  $N(y) \subseteq \{x, z, u, u', w, w', w''\}$ . Since  $y \notin V_5$ , we see that  $|N(y) \cap \{w, w', w''\}| \ge 2$ . Without loss of generality, we may assume that  $\{w, w'\} \subseteq N(y)$ . Let B' be a fragment with respect to zu and let T' = N(B'). Since  $N(z) \cap N(u) \subseteq \{y\}$  and  $y \notin V_5$ , we see that  $N(z) \cap N(u) \cap V_5 = \emptyset$ , which implies that neither B' nor  $\overline{B'}$  is trivial, and hence  $|B'| \ge 2$  and  $|\overline{B'}| \ge 2$ . Since  $S' - \{z\} \subseteq N(u')$ , we see that  $u' \in T'$ .

### Claim 8.1 $y \in T'$ .

*Proof.* Assume  $y \notin T'$ . Without loss of generality, we may assume that  $y \in B'$ . Then, since  $\{w, w'\} \subseteq N(y)$ ,  $\{w, w'\} \subseteq T' \cup B'$ . Hence, we observe that  $N(\{u, u'\}) \cap \overline{B}' = \{w''\}$ . Then, assures us that  $\overline{B}' = \{w''\}$ , which contradicts the previous observation that  $|\overline{B}'| \ge 2$ . This contradiction proves Claim 8.1.

By Claim 8.1, we see that  $\{y, z, u, u'\} \subseteq T'$ , which implies  $N(u) \cap (B' \cup \overline{B}') \subseteq \{w, w', w''\}$ since  $N(u) \subseteq \{y, z, u', w, w', w''\}$ . We also observe that  $N(u') \cap (B' \cup \overline{B}') \subseteq \{w, w', w''\}$  since  $N(u') = \{u, y, w, w', w''\}$ . Since neither  $N(u) \cap B' = \emptyset$  nor  $N(u) \cap \overline{B}' = \emptyset$ , we have either  $|B' \cap \{w, w', w''\}| = 1$  or  $|\overline{B}' \cap \{w, w', w''\}| = 1$ . Without loss of generality, we may assume that  $|B' \cap \{w, w', w''\}| = 1$ , say  $B' \cap \{w, w', w''\} = \{\tilde{w}\}$ . Then we see that  $N(\{u, u'\}) \cap B' = \{\tilde{w}\}$  and applying Lemma 1 with the roles of A and S replaced by B' and  $\{u, u'\}$ , respectively, we see that  $B' = \{\tilde{w}\}$ , which contradicts the previous observation that  $|B'| \ge 2$ . This contradiction proves that z is a strongly admissible vertex for (x, y; A) and (1) is shown.

Next we show (2). Assume z is not a hyper admissible vertex for (x; A). We show  $(N(z) \cap S - \{x\}) \cap (V_5 - V_5^{(2)}) \neq \emptyset$ . Since z is strongly admissible and not hyper admissible, we see that  $|N(z) \cap A| = 2$ ,  $|N(z) \cap S| = 2$ ,  $|N(z) \cap \overline{A}| = 1$  and  $|S \cap \overline{B}| = 2$ . Let  $N(z) \cap A = \{y, u\}$ ,  $N(z) \cap S = \{x, w\}$ ,  $N(z) \cap \overline{A} = \{v\}$  and  $S \cap \overline{B} = \{w', w''\}$ . Let  $A' = A - \{y\}$  and  $S' = N(A') = (S - \{x\}) \cup \{y\} = \{z, y, w, w', w''\}$ . Then A' is a fragment of G such that  $|A'| \ge 2$  and  $|\overline{A}'| = |\overline{A} \cup \{x\}| \ge 3$ . Note that  $N(z) \cap A' = \{u\}$ .

### Claim 8.2 w is an admissible vertex for (z, u; A').

*Proof.* At first we consider the case that  $|A'| \ge 3$ . In this case we have  $|A'| \ge 3$ ,  $|\bar{A}'| \ge 3$ and  $N(z) \cap A' = \{u\}$ . Thus Lemma 3 assures us the existence of an admissible vertex for (z, u; A'). Since  $N(z) \cap S' = \{y, w\}$  and  $y \notin V_5$ , we observe that w is an admissible vertex for (z, u; A').

Next we consider the case that |A'| = 2, say  $A' = \{u, u'\}$ . Since  $A' \cap B = \emptyset$  and  $A' \cap T = \{u\}$ , we see that  $u' \in A' \cap \overline{B}$  and  $N(u') = \{y, u, w, w', w''\}$ . Since  $N(y) \subseteq S \cup A$  and  $A = \{y, u, u'\}$ , the fact  $y \in V_{\geq 6}$  implies  $|N(y) \cap \{w, w', w''\}| \geq 2$ . Let B' be a fragment with respect to zu and let T' = N(B'). Since  $S' - \{z\} \subseteq N(u')$ , we observe that  $u' \in T'$ , which implies  $A' \cap T' = \{u, u'\}$  and  $A' \cap B' = A' \cap \overline{B}' = \emptyset$ . Since  $A' \cap B' = A' \cap \overline{B}' = \emptyset$ , we see that neither  $S' \cap B' = \emptyset$  nor  $S' \cap \overline{B}' = \emptyset$ . We show that either  $|S' \cap B'| = 1$  or  $|S' \cap \overline{B}'| = 1$ . If  $y \in S' \cap T'$ , then  $|S' \cap T'| \ge 2$ , which implies either  $|S' \cap B'| = 1$  or  $|S' \cap \overline{B}'| = 1$ . Hence assume  $y \notin S' \cap T'$ . If  $y \in S' \cap \overline{B}'$ , then, the fact that  $|N(y) \cap \{w, w', w''\}| \ge 2$  assures us that  $|S' \cap B'| = 1$ . Similarly, if  $y \in S' \cap B'$ , then we have  $|S' \cap \overline{B}'| = 1$ . Now it is shown that either  $|S' \cap B'| = 1$  or  $|S' \cap B'| = 1$  or  $|S' \cap B'| = 1$ .

Without loss of generality, we may assume that  $|S' \cap B'| = 1$ , say  $S' \cap B' = \{\tilde{w}\}$ . Then, since  $|S' \cap B'| < |A' \cap T'|$ , we observe that  $\bar{A}' \cap B' = \emptyset$  and  $B' = S' \cap B' = \{\tilde{w}\}$ . Hence we know that  $\tilde{w} \in V_5$  and  $\tilde{w}z \in E(G)$ . Since  $N(z) \cap S' = \{y, w\}$  and  $y \notin V_5$ , we see that  $\tilde{w} = w$ , which implies the desired conclusion that w is an admissible vertex for (z, u; A').

If  $w \notin V_5^{(2)}$ , then  $w \in (N(y) \cap S - \{x\}) \cap (V_5 - V_5^{(2)})$  and we are done. Hence assume  $w \in V_5^{(2)}$ .

Claim 8.3 If  $N(w) \cap \overline{A} \cap V_5 = \emptyset$ , then  $|\overline{A}| \ge 3$ .

*Proof.* Assume  $N(w) \cap \overline{A} \cap V_5 = \emptyset$ . Since  $\overline{A} \cap B = \emptyset$ ,  $\overline{A} \cap T = \{v\}$  and  $|\overline{A}| \ge 2$ , we observe  $\overline{A} \cap \overline{B} \neq \emptyset$ , which implies  $\overline{A} \cap \overline{B}$  is a fragment of G since  $|(S \cap \overline{B}) \cup (S \cap T) \cup (\overline{A} \cap T)| = 5$ . Hence  $N(w) \cap (\overline{A} \cap \overline{B}) \neq \emptyset$ , say  $v' \in N(w) \cap (\overline{A} \cap \overline{B})$ . Then, since  $N(w) \cap \overline{A} \cap V_5 = \emptyset$ , we see that  $v' \notin V_5$ , which implies  $|\overline{A} \cap \overline{B}| \ge 2$ . This implies the desired conclusion  $|\overline{A}| = |\overline{A} \cap T| + |\overline{A} \cap \overline{B}| \ge 3$ .

**Claim 8.4**  $N(w) \cap A \cap V_5 = \emptyset$ .

**Proof.** Assume  $N(w) \cap A \cap V_5 \neq \emptyset$ . Then, since  $z \in N(w) \cap V_5$  and  $w \in V_5^{(2)}$ , we see that  $N(w) \cap \overline{A} \cap V_5 = \emptyset$ . Hence Claim 8.3 assures us that  $|\overline{A}| \ge 3$ . Since  $|\overline{A}|, |A| \ge 3$  and  $N(w) \cap \overline{A} \cap V_5 = \emptyset$ , applying Corollary 7, we see that there is an admissible vertex for  $(w; \overline{A})$ . Since  $z \in N(w) \cap V_5$ ,  $N(w) \cap A \cap V_5 \neq \emptyset$  and  $w \in V_5^{(2)}$ , we observe that  $N(w) \cap S \cap V_5 = \{z\}$ . Since  $|N(z) \cap \overline{A}| = 1$ , z is not an admissible vertex for  $(w; \overline{A})$ , which implies that there is no admissible vertex for  $(w; \overline{A})$ . This contradicts the previous assertion and this contradiction proves Claim 8.4.

**Claim 8.5**  $|A'| \ge 3$ .

*Proof.* Since  $A' \cap B = \emptyset$ ,  $A' \cap T = \{u\}$  and  $|A'| \ge 2$ , we observe that  $A' \cap \overline{B} \ne \emptyset$ , which implies that  $A' \cap \overline{B}$  is a fragment of G since  $|(S \cap \overline{B}) \cup (S \cap T) \cup (A \cap T)| = 5$ , which implies  $|A' \cap \overline{B}| \ge 2$ . This implies the desired conclusion that  $|A'| = |A' \cap T| + |A' \cap \overline{B}| \ge 3$ .

We proceed with the proof of Lemma 8 (2).

Since  $|\bar{A}'|, |A'| \ge 3$  and  $N(w) \cap A' \cap V_5 = \emptyset$ , applying Corollary 7, we see that there is an admissible vertex  $\tilde{w}$  for (w; A'). Since  $|N(z) \cap A'| = 1$ , z is not an admissible vertex for (w; A'), which implies  $\tilde{w} \ne z$ . Then, since  $w \in V_5^{(2)}$ , we observe that  $N(w) \cap V_5 = \{z, \tilde{w}\}$ , which implies that  $N(w) \cap \bar{A}' \cap V_5 = \emptyset$ . Since  $\bar{A} = \bar{A}' - \{x\}$ ,  $N(w) \cap \bar{A}' \cap V_5 = \emptyset$  implies  $N(w) \cap \bar{A} \cap V_5 = \emptyset$ . Now we have  $N(w) \cap \bar{A} \cap V_5 = \emptyset$  and Claim 8.3 assures us that  $|\bar{A}| \ge 3$ . Since  $|\bar{A}|, |A| \ge 3$ ,  $|N(w) \cap \bar{A}| = 1$ 

and  $N(w) \cap \overline{A} \cap V_5 = \emptyset$ , applying (1), we see that there is a strongly admissible vertex for  $(w; \overline{A})$ . However, since  $N(w) \cap S \cap V_5 = \{z, \tilde{w}\}$ ,  $|N(z) \cap A| \ge 2$  and  $|N(\tilde{w}) \cap A| \ge 2$ , we see that there is no strongly admissible vertex for  $(w; \overline{A})$ , which violates the previous assertion. This contradiction proves (2) and the proof of Lemma 8 is completed.

# 4 The proof of Theorem 1

In this section we give a proof of Theorem 1.

Let G be a 5-connected graph. Let A be a fragment of G and let S = N(A). Let Ad (Y;A) denote the set of admissible vertices for (Y;A). We demote  $\hat{S}_A$  the set of vertices y of S such that  $Ad(y;A) \neq \emptyset$  and let  $\tilde{S}_A = \bigcup_{y \in \hat{S}_A} Ad(y;A)$ . Using these notation, we can rewrite Theorem 1 as the following.

**Theorem 1** Let G be a contraction-critically 5-connected graph. Let A be a connected fragment of G with |A| = 2, say  $A = \{x_1, x_2\}$  and let S = N(A).

- (1) If  $A \cap V_6 \neq \emptyset$ , then  $|\hat{S}_A| \ge 4$ .
- (2) If  $A \cap V_6 \neq \emptyset$ ,  $|\tilde{S}_A| \ge 3$ .
- (3) If  $A \cap V_6 = \emptyset$ , then  $\hat{S}_A \cap (S N(x_1) \cap N(x_2)) \neq \emptyset$ .

We prove Theorem 1 using the notation  $\hat{S}_A$  and  $\tilde{S}_A$ . Let  $S = \{y_1, y_2, y_3, y_4, y_5\}$ . Without loss of generality we may assume that  $\deg_G(x_1) \ge \deg_G(x_2)$ . Hence, if  $A \cap V_6 \ne \emptyset$ , then  $x_1 \in V_6$  and  $S \subseteq N(x_1)$ .

(1) Assume  $A \cap V_6 \neq \emptyset$  and  $|\hat{S}_A| \leq 3$ . Then  $|S - \hat{S}_A| \geq 2$ , say  $y_1, y_2 \in S - \hat{S}_A$ .

We show that there is a fragment  $B_i$  such that  $\{x_1, x_2, y_i\} \subseteq N(B_i)$  for i = 1, 2. Let  $i \in \{1, 2\}$ . If  $x_2y_i \in E(G)$ , then let  $B_i$  be a fragment with respect to  $x_2y_i$ . Then, since  $S \subseteq N(x_1)$ , we observe that  $\{x_1, x_2, y_i\} \subseteq N(B_i)$ . If  $x_2y_i \notin E(G)$ , then let  $B_i$  be a fragment with respect to  $x_1y_i$ . Then, since  $S - \{y_i\} \subseteq N(x_2)$ , we again observe that  $\{x_1, x_2, y_i\} \subseteq N(B_i)$ . Now the existence of a fragment  $B_i$  such that  $\{x_1, x_2, y_i\} \subseteq N(B_i)$  is shown.

Let  $B_i$  be a fragment such that  $\{x_1, x_2, y_i\} \subseteq N(B_i)$  and let  $T_i = N(B_i)$  for i = 1, 2. We show that  $|S \cap B_1| \ge 2$ . Suppose  $|S \cap B_1| \le 1$ . Then, since  $|S \cap B_1| < |A \cap T_1|$ , Lemma 2 (2) assures us that  $\overline{A} \cap B_1 = \emptyset$ , which implies  $B_1 = S \cap B_1$ , say  $B_1 = S \cap B_1 = \{y\}$ . Then we observe  $y \in V_5$ and  $\{y_1\} \cup A \subseteq N(y)$ , which implies that y is an admissible vertex for  $(y_1; A)$ . This contradicts the fact that  $y_1 \in S - \hat{S}_A$  and it is shown that  $|S \cap B_1| \ge 2$ .

By the similar arguments, we can show that  $|S \cap \overline{B}_1|, |S \cap B_2|, |S \cap \overline{B}_2| \ge 2$ . Thus we have

 $|S \cap B_i| = |S \cap \overline{B}_i| = 2$  for i = 1, 2. Without loss of generality we may assume that  $S \cap B_1 = \{y_2, y_3\}$  and  $S \cap \overline{B}_1 = \{y_4, y_5\}$ . Say  $S \cap B_2 = \{y_1, y_j\}$  and  $S \cap \overline{B}_2 = \{y_3, y_4, y_5\} - \{y_j\}$ . Then we observe that  $y_1 \in T_1 \cap B_2$  and  $y_2 \in T_2 \cap B_1$ .

We show  $j \neq 3$ . Suppose j = 3. Then  $y_3 \in B_1 \cap B_2$  and  $y_4, y_5 \in \overline{B}_1 \cap \overline{B}_2$ . Since neither  $B_1 \cap B_2$  nor  $\overline{B}_1 \cap \overline{B}_2$  is empty, we see that  $B_1 \cap B_2$  is a fragment of G. Since  $\{x_1, x_2\} \subseteq N(B_1 \cap B_2)$  and  $N(\{x_1, x_2\}) \cap (B_1 \cap B_2) = \{y_3\}$ , applying Lemma 1 with the roles of A and S replaced by  $B_1 \cap B_2$  and  $\{x_1, x_2\}$ , respectively, we see that  $B_1 \cap B_2 = \{y_3\}$ , which implies  $y_3 \in V_5$  and  $\{y_1\} \cup A \subseteq N(y_3)$ . Hence  $y_3 \in Ad(y_1; A)$ , which contradicts the choice of  $y_1$ . This contradiction shows  $j \neq 3$ , say j = 4.

In this situation, we observe that  $y_3 \in B_1 \cap \overline{B}_2$ ,  $y_4 \in \overline{B}_1 \cap B_2$  and  $y_5 \in \overline{B}_1 \cap \overline{B}_2$ . Since neither  $\overline{B}_1 \cap B_2$  nor  $B_1 \cap \overline{B}_2$  is empty, we see that  $\overline{B}_1 \cap B_2$  is a fragment of G. Since  $\{x_1, x_2\} \subseteq N(\overline{B}_1 \cap B_2)$  and  $N(\{x_1, x_2\}) \cap (\overline{B}_1 \cap B_2) = \{y_4\}$ , applying Lemma 1 with the roles of A and S replaced by  $\overline{B}_1 \cap B_2$  and  $\{x_1, x_2\}$ , respectively, we see that  $\overline{B}_1 \cap B_2 = \{y_4\}$ , which implies  $y_4 \in V_5$  and  $\{y_1\} \cup A \subseteq N(y_4)$ . Hence  $y_4 \in Ad(y_1; A)$ , which again contradicts the choice of  $y_1$ . This contradiction shows that  $|\hat{S}_A| \ge 4$  and (1) is proved.

(2) Assume  $A \cap V_6 \neq \emptyset$  and  $|\tilde{S}_A| \leq 2$ . Since  $\hat{S}_A \neq \emptyset \geq 4$ , we see that  $\tilde{S}_A \neq \emptyset$ , say  $y \in \tilde{S}_A$ . Since  $y \in V_5$ ,  $A \subseteq N(y)$  and  $N(y) \cap \bar{A} \neq \emptyset$ , we see that  $|N(y) \cap S| \leq 2$ . Since  $|\hat{S}_A| \geq 4$  and  $|N(y) \cap S| \leq 2$  for  $y \in \tilde{S}_A$ , we see that  $|\tilde{S}_A| = 2$ , say  $\tilde{S}_A = \{y_1, y_2\}$ . Since  $|\hat{S}_A| \geq 4$ , we see that either  $y_1 \in \hat{S}_A$  or  $y_2 \in \hat{S}_A$ , which implies  $y_1y_2 \in E(G)$  and  $\{y_1, y_2\} \subseteq \hat{S}_A$ . Since  $|N(y) \cap S| \leq 2$  for  $y \in \tilde{S}_A$  and  $y_1y_2 \in E(G)$ , we see that  $|N(\tilde{S}_A) \cap S| \leq 2$ , which implies  $|\hat{S}_A| \leq 4$  and  $S - \hat{S}_A \neq \emptyset$ , say  $y_j \in S - \hat{S}_A$ .

By the same arguments in the proof of (1), we see there is a fragment B with  $\{x_1, x_2, y_j\} \subseteq N(B)$  and we also have  $|S \cap B| = |S \cap \overline{B}| = 2$ . Let T = N(B). Since  $y_1y_2 \in E(G)$ , we may assume that  $S \cap B = \{y_1, y_2\}$ . Since  $E_G(S \cap B, S \cap \overline{B}) = \emptyset$  and  $S \cap B = \tilde{S}_A$ , we see that  $N(\tilde{S}_A) \cap (S \cap \overline{B}) = \emptyset$ , which implies  $\hat{S}_A \cap (S \cap \overline{B}) = \emptyset$  and  $|\hat{S}_A| = |\tilde{S}_A| = 2$ . This contradicts (1) and this contradiction shows  $|\tilde{S}_A| \geq 3$ . Now (2) is proved.

(3) Assume  $\hat{S}_A \cap (S - N(x_1) \cap N(x_2)) = \emptyset$ . Since |A| = 2 and  $A \cap V_6 = \emptyset$ , we observe that  $A \subseteq V_5$ , which implies  $|N(x_1) \cap N(x_2)| = 3$ , say  $N(x_1) \cap N(x_2) = \{y_3, y_4, y_5\}$  and  $N(x_i) - \{x_{3-i}, y_3, y_4, y_5\} = \{y_i\}$  for i = 1, 2. Then  $S - N(x_1) \cap N(x_2) = \{y_1, y_2\}$  and by the assumption, we observe that  $\hat{S}_A \cap \{y_1, y_2\} = \emptyset$ . Since  $\hat{S}_A \cap \{y_1, y_2\} = \emptyset$ , we observe that  $Ad(y_i; A) = \emptyset$  for i = 1, 2. Let  $B_i$  be a fragment with respect to  $x_i y_i$  and let  $T_i = N(B_i)$  for i = 1, 2.

We show that  $x_2 \in T_1$ . Suppose  $x_2 \notin T_1$ , say  $x_2 \in \overline{B}_1$ . Then since  $N(x_2) = A \cup S - \{x_2, y_1\}$ , we observe that  $A \cup S \subseteq \overline{B}_1 \cup T_1$ , which implies  $N(x_1) \cap B_1 = \emptyset$ . This contradicts the choice of  $B_1$ and it is shown that  $x_2 \in T_1$ . Similarly we have  $x_1 \in T_2$ . Now we know that  $\{x_1, x_2, y_1, y_2\} \subseteq T_1 \cup T_2$ . Since  $\{x_1, x_2, y_1, y_2\} \subseteq T_1 \cup T_2$  and neither  $N(x_1) \cap B_1$  nor  $N(x_1) \cap \bar{B}_1$  is empty, we see neither  $B_1 \cap \{y_3, y_4, y_5\}$  nor  $\bar{B}_1 \cap \{y_3, y_4, y_5\}$  is empty, which implies either  $|B_1 \cap \{y_3, y_4, y_5\}| = 1$  or  $|\bar{B}_1 \cap \{y_3, y_4, y_5\}| = 1$ , say  $|B_1 \cap \{y_3, y_4, y_5\}| = 1$  and let  $B_1 \cap \{y_3, y_4, y_5\} = \{y_3\}$ .

We show that  $y_2 \in B_1$ . Suppose  $y_2 \notin B_1$ . Then, since  $N(\{x_1, x_2\}) \cap B_1 = \{y_3\}$ , applying Lemma 1 with the roles of A and S replaced by  $B_1$  and  $\{x_1, x_2\}$ , respectively, we see that  $B_1 = \{y_3\}$ , which implies  $y_3 \in V_5$  and  $\{y_1\} \cup A \subseteq N(y_3)$ . Hence we have  $y_3 \in Ad(y_1; B_1)$ , which contradicts the assumption that  $\hat{S}_A \cap \{y_1, y_2\} = \emptyset$ . This contradiction shows that  $y_2 \in B_1$ .

We show that  $\{y_4, y_5\} \subseteq \overline{B}_1$ . Suppose  $y_5 \notin \overline{B}_1$ . Then, since  $N(\{x_1, x_2\}) \cap \overline{B}_1 = \{y_4\}$ , applying Lemma 1 with the roles of A and S replaced by  $\overline{B}_1$  and  $\{x_1, x_2\}$ , respectively, we see that  $\overline{B}_1 = \{y_4\}$ , which implies  $y_4 \in Ad(y_1; A)$ . This contradicts the assumption that  $\hat{S}_A \cap \{y_1, y_2\} = \emptyset$ and it is shown that  $y_5 \in \overline{B}_1$ .

By symmetry we have  $y_4 \in \overline{B}_1$ . By the similar argument, we know that there is an integer  $j \in \{3, 4, 5\}$  such that  $\{y_1, y_j\} \subseteq B_2$  and  $\{y_3, y_4, y_5\} - \{y_j\} \subseteq \overline{B}_2$ . In this situation we observe that  $y_1 \in T_1 \cap B_2$  and  $y_2 \in T_2 \cap B_1$ .

We show  $j \neq 3$ . Suppose j = 3. Then  $y_3 \in B_1 \cap B_2$  and  $y_4, y_5 \in \overline{B}_1 \cap \overline{B}_2$ . Since neither  $B_1 \cap B_2$  nor  $\overline{B}_1 \cap \overline{B}_2$  is empty, Lemma 2 (1) assures us that  $B_1 \cap B_2$  is a fragment of G. Since  $\{x_1, x_2\} \subseteq N(B_1 \cap B_2)$  and  $N(\{x_1, x_2\}) \cap (B_1 \cap B_2) = \{y_3\}$ , applying Lemma 1, with the roles of A and S replaced by  $B_1 \cap B_2$  and  $\{x_1, x_2\}$ , respectively, we see that  $B_1 \cap B_2 = \{y_3\}$ , which implies  $y_3 \in V_5$  and  $\{y_1\} \cup A \subseteq N(y_3)$ . Hence  $y_3 \in Ad(y_1; A)$ , which contradicts the assumption that  $\hat{S}_A \cap \{y_1, y_2\} = \emptyset$  and it is shown that  $j \neq 3$ , say j = 4.

Then  $y_3 \in B_1 \cap \bar{B}_2$ ,  $y_4 \in \bar{B}_1 \cap B_2$  and  $y_5 \in \bar{B}_1 \cap \bar{B}_2$ . Since neither  $\bar{B}_1 \cap B_2$  nor  $B_1 \cap \bar{B}_2$  is empty, we see that  $\bar{B}_1 \cap B_2$  is a fragment of G. Since  $\{x_1, x_2\} \subseteq N(\bar{B}_1 \cap B_2)$  and  $N(\{x_1, x_2\}) \cap (\bar{B}_1 \cap B_2) = \{y_4\}$ , applying Lemma 1 with the roles of A and S replaced by  $\bar{B}_1 \cap B_2$  and  $\{x_1, x_2\}$ , respectively, we see that  $\bar{B}_1 \cap B_2 = \{y_4\}$ , which implies  $y_4 \in V_5$  and  $\{y_1\} \cup A \subseteq N(y_4)$ . Hence  $y_4 \in Ad(y_1; A)$ , which contradicts the assumption that  $\hat{S}_A \cap \{y_1, y_2\} = \emptyset$ . This contradiction shows that  $\hat{S}_A \cap (S - N(x_1) \cap N(x_2)) \neq \emptyset$ . Now (3) is proved and the proof of Theorem 1 is completed.

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